Enhancing Performance And Reliability of Rule Management Platforms

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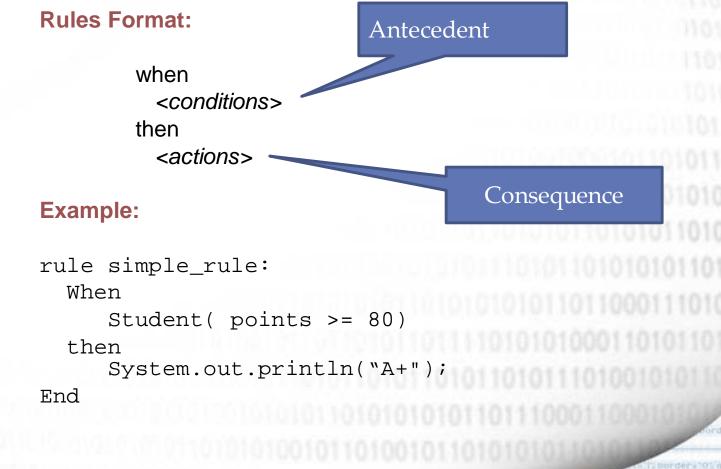
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RulE Management Platforms (REMPs) allow software engineers to represent programming logic as conditional sentences that relate statements of facts (i.e., **rules**) using high-level declarative languages



Rule Structure And Example



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An Interesting Question

Rule A:

When Product (Price < 80) Then set(Price, 100);

Rule B:

When Product (Price > 50) Then set(Price, 200);

Price = 65

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If-Then Decision Points

Alternative to rules is the imperative model where sequences of ifthen statements with conditionals and loops are evaluated in a strictly defined order.

A weird program for calculating Pi written in Fortran. 1 C From: Fink, D.G., Computers and the Human Mind, Anchor Books, 1966. 2 C PROGRAM PI DIMENSION TERM(100) N=1 TERM(N)=((-1)**(N+1))*(4./(2.*N-1.)) N=N+1 IF (N-101) 3,6,6 N=1SUM98 = SUM98 + TERM(N)12 WRITE(*,28) N, TERM(N) N=N+1IF (N-99) 7, 11, 11 SUM99=SUM98+TERM(N) SUM100=SUM99+TERM(N+1) IF (SUM98-3.141592) 14,23,23 17 IF (SUM99-3.141592) 23,23,15 IF (SUM100-3.141592) 16,23,23 G16 AV89=(SUM98+SUM99)/2. 20 AV90=(SUM99+SUM100)/2. 21 22 COMANS=(AV89+AV90)/2. IF (COMANS-3.1415920) 21,19,19 23 IF (COMANS-3.1415930) 20,21,21 20 WRITE(*,26) 25 GO TO 22 26 27 WRITE(*,27) COMANS **STOP** 28 22 29 23 WRITE(*,25) GO TO 22 FORMAT ('ERROR IN MAGNITUDE OF SUM') 26 FORMAT('PROBLEM SOLVED') 33 27 FORMAT('PROBLEM UNSOLVED', F14.6) FORMAT(I3, F14.6) 34 28 35 END

If-Then Decision Points

A weird program for calculating Pi written in Fortran. 1 C Alternative to rules From: Fink, D.G., Computers and the Human Mind, Anchor Books, 1966. PROGRAM PI DIMENSION TERM(100) is the Hard to maintain and mod inefficient code that is sequ statemot adaptable to then with frequent changes in and business requirements eval FORMAT('PROBLEM SOLVED') strictly defined 27 FORMAT('PROBLEM UNSOLVED', F14.6) 33 FORMAT(I3, F14.6) 28 35 END order.

REMPS Are Widely Used

According to market report from Forrester, the estimated revenue of business rules management systems (BRMS) increased from \$265 million in 2008 to over \$600 million in 2011. It is one of the fastest growing markets.

One leading vendor, IBM ILOG Optimization is used by over 50% of the world's largest companies, 1000's of Universities, and 1000's of application providers.

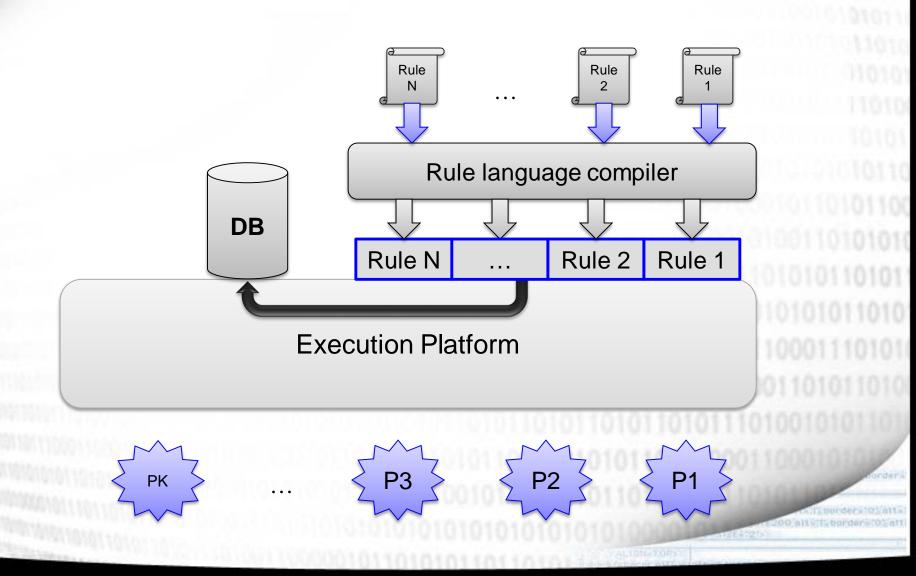


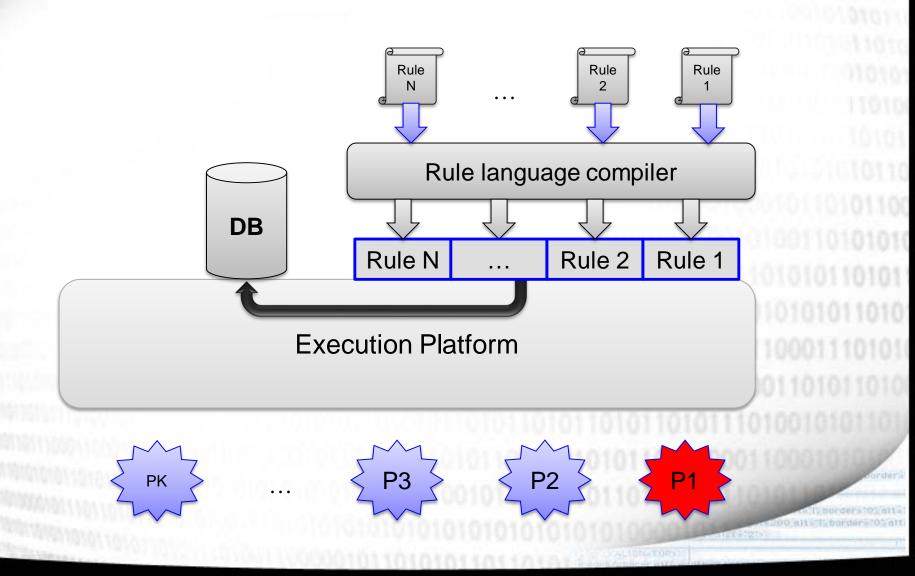
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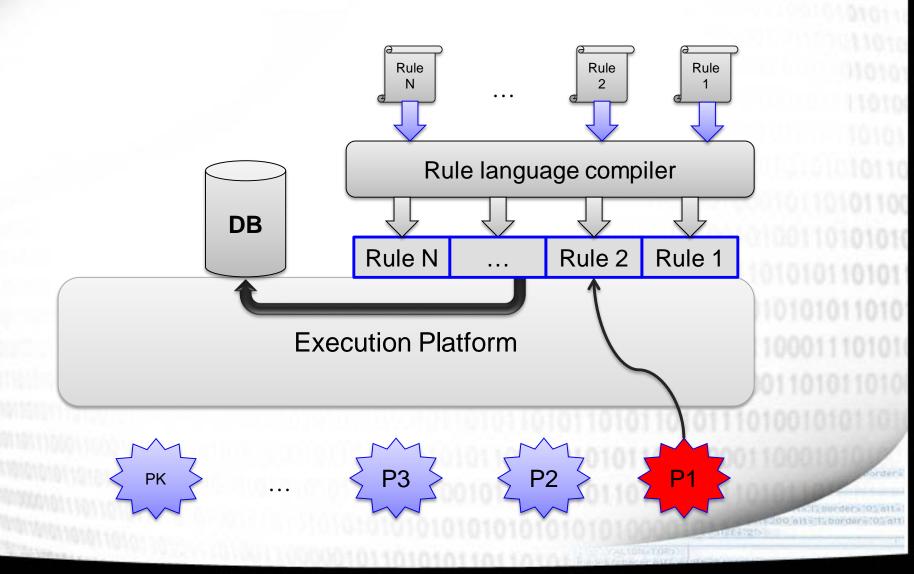


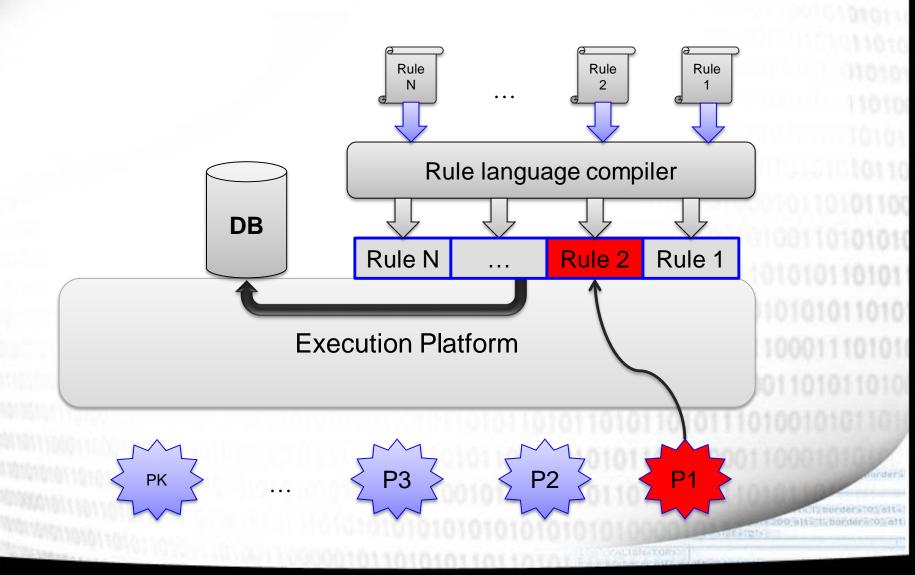
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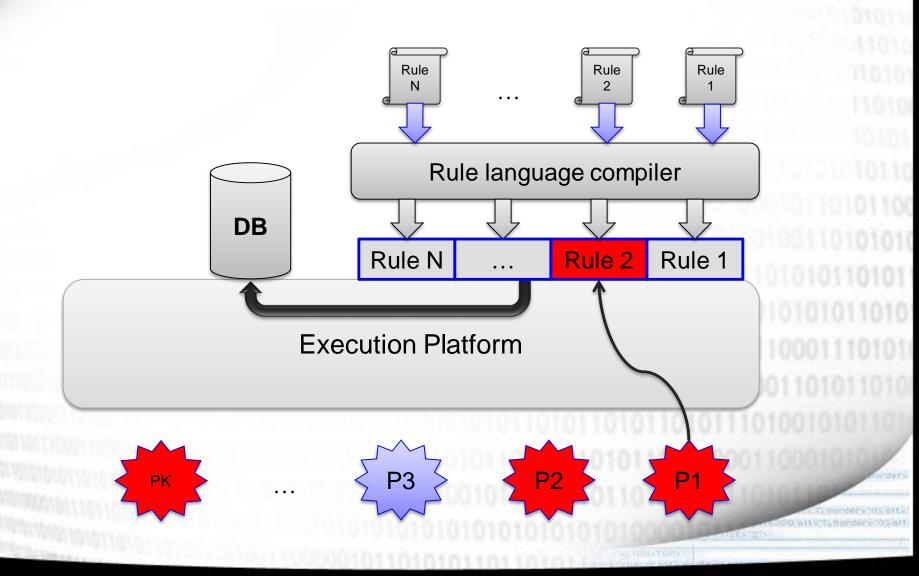


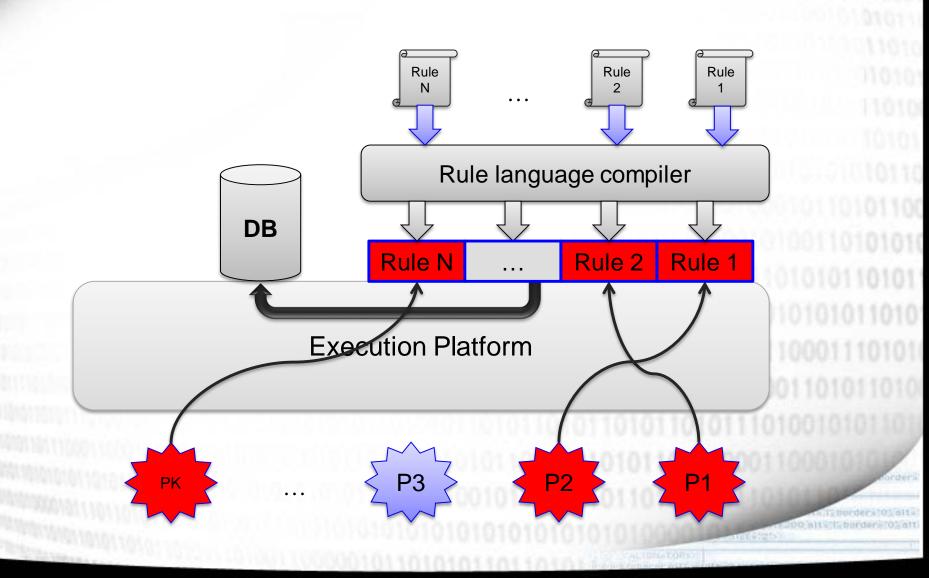


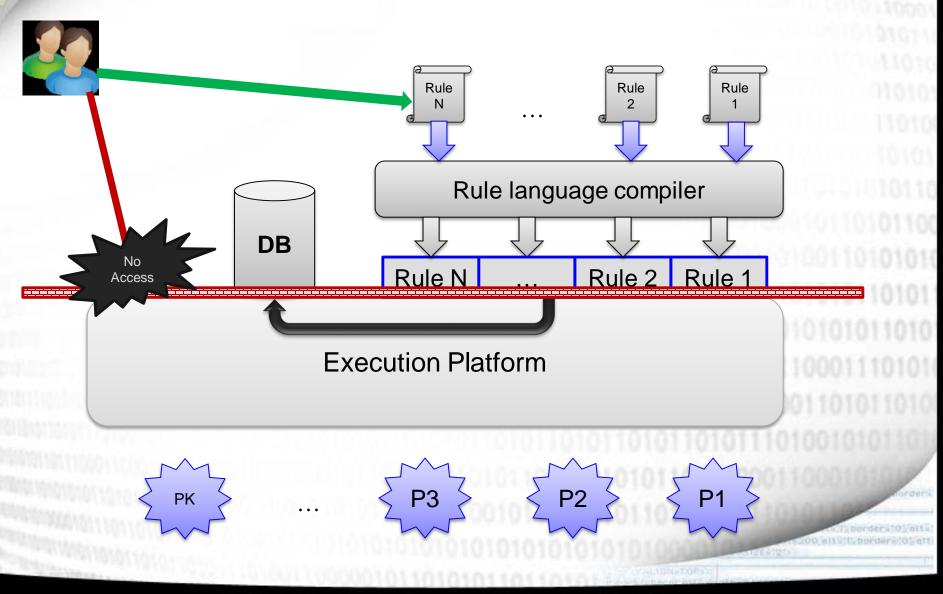












Separation of Concerns

A key property of REMPs is that they encapsulate the control flow that includes fact inference and rule firing logics

 It is a fundamental separation of concerns of the control flow and the rule business logic.

Software engineers concentrate on reasoning about higherlevel business logic that they encode in rules without worrying about low-level details of rule invocations by effectively delegating this job to REMP engines

 Rule-driven APplications (RAPs) are highly adaptable to changing requirements, since stakeholders simply add new rules as independent modules to RAPs

Benefits of REMPs And RAPs

Easy to maintain & evolve

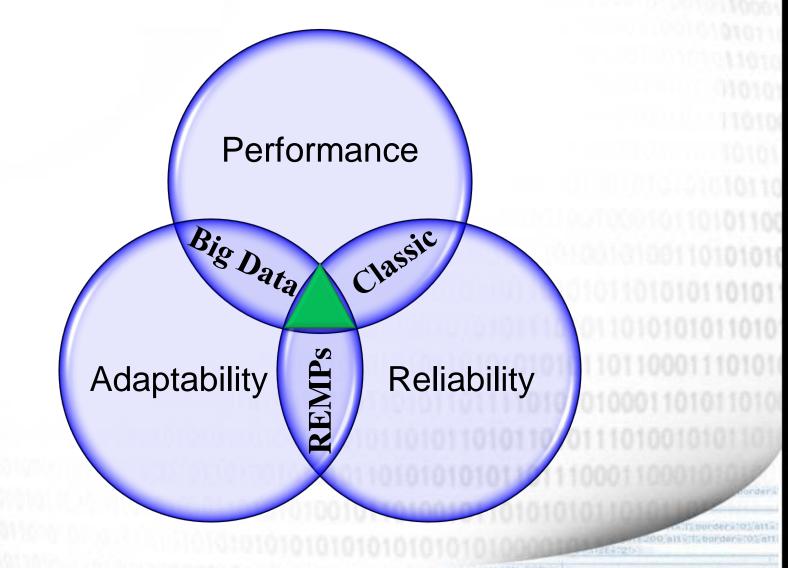
Rules are easy to comprehend and highly modular

RAPs are highly adaptable

Constraints of REMPs And RAPs								
RAPs may contain tens of thousands of rules								
Detection of conflicting rules is difficult								
Dependencies should not be introduced among rules	0101000110101010101010101010101010101010							

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The PAR Model For REMPs



Fundamental Problem

Many REMPs execute RAPs sequentially.

Locks introduce complex dependencies among rules, thereby defeating the separation of concerns and eventually the adaptability of RAPs.

How to enhance the performance of RAPs without sacrificing their adaptability and reliability?

Motivating Example: Credit

rule "Rule-Credit" salience 10 when

\$cashflow : Cashflow(\$account:account,

\$date : date, \$amount : amount,

type==Cashflow.CREDIT)

not Cashflow(account==\$account,date<\$date)</pre>

then

//some code

\$account.setBalance(

\$account.getBalance()+\$amount);

retract(\$cashflow);

end

Motivating Example: Debit

rule "Rule-Debit" salience 1 when

\$cashflow : Cashflow(\$account : account,

\$date : date, \$amount : amount,

type==Cashflow.DEBIT)

not Cashflow(account==\$account,date<\$date) then

//some code

if(\$account.getBalance()>\$amount){
 \$account.setBalance(

\$account.getBalance()-\$amount); }
else { new BlockedAccount(\$cashflow); }
retract(\$cashflow);

end

Parallelism Interferes With Saliences

Let us assume that a REMP engine executes rules in parallel and lock objects are used to synchronize concurrent accesses.

Using a lock object effectively overrides the intention of the programmer to give the priority to the rule with a higher salience.

Given the large number of possible interleavings among tens of thousands of rules in a RAP, it is very difficult to reason about interactions between saliences and synchronization lock mechanisms.

Reliability Meets Performance

In fully parallelized REMPs, loss or reliability comes from two sources:

- different orders in which rules are executed by the REMP engines;
- races between parallelized executions of rules.

When the system produces different results consecutively for the same computational task using the same input data, it is a serious problem, since it reduces the confidence of the user in the RAP and it impacts negatively the perception of the user about the business value that the company or organization delivers.

- the execution order for different instructions can be affected by multiple factors beyond the control of stakeholders
- sometimes even slight changes in the non-functional parameters of the environment (e.g., paging on demand) for executing RAPs result in different orders of instruction interleavings that lead to different results, hence the loss of reliability

The Problem Statement

Enable REMPs to execute rules in RAPs in parallel

Do not violate the separation of concerns in REMPs by requiring programmers to use synchronization lock mechanisms for concurrent accesses to shared resources

Prevent races in parallelized RAPs without explicit using of locking mechanisms by programmers

Choose a better schedule for executing rules that share the same lock objects to improve the overall performance of RAPs

Core Ideas



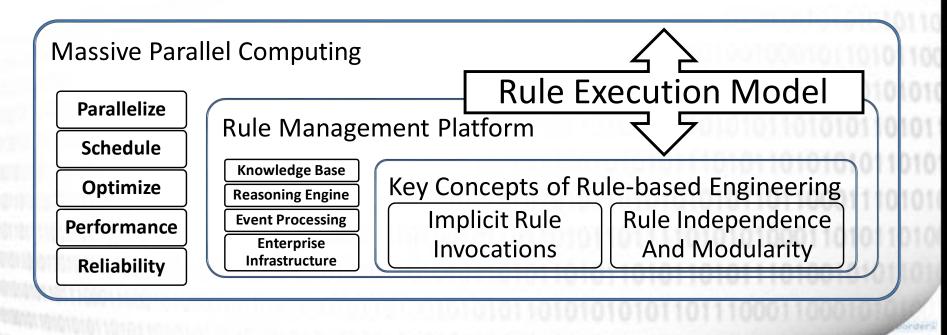
- Find all concurrent access to resources from rules where one of the accesses is write.
- 2) Define synchronizations around these accesses.
 - 3) Impose a complete ordering among all rules that are fired in working memory.

Our Solution - PERLATO

PErformance and Reliability for ruLe-driven ApplicaTiOns (PERLATO) connects separate layers or REMPs in a way that enable us to solve the fundamental problem of REMP.

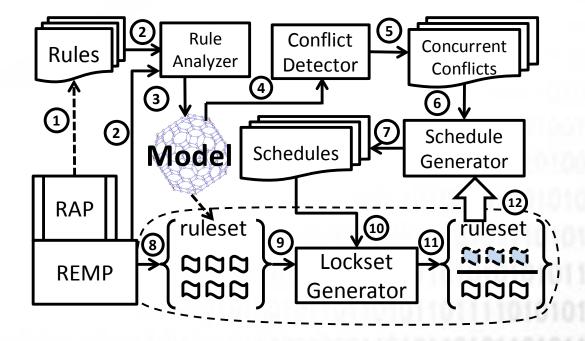
- we obtain a rule execution model from a RAP that approximate different execution scenarios by using the if-then structure of rules by analyzing their antecedents and consequents
- the obtained rule execution model is used in PERLATO to detect races statically among these rules effectively and efficiently
- the rule execution model and locking strategies for a given RAP are passed to the REMP engine, so that it can precompute an execution schedule for rules in a RAP to optimize the performance of the RAP.
- We implemented PERLATO for JBoss Drools, an open-source enterprise-level REMP and we evaluated PERLATO on three RAPs. The results suggest that PERLATO is effective and efficient, since we achieved up to 225% speedup on average without observing any races.

Map fo PERLATO



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The Architecture of PERLATO



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Research Questions

Is PERLATO effective in achieving higher speedups for subject RAPs?

Is finer granularity locking strategy more effective in obtaining higher speedup for RAPs?

Is symbiotic scheduling effective in obtaining higher speedup for RAPs?

Result for Subject RAPs

	Input Set	Facts	Rules Fired	Approach	Type	Avg	Med	Min	Max	σ^2	100010001
EEWS		246896	1088	DROOLS	Seq	32.22	32.08	31.559	35.358	0.4	101010111
	1				Par	9.52	9.27	8.267	11.789	0.87	11111111111
				PERLATO	Rule	13.44	13.6	10.927	14.911	0.99	1.11010
					Atomic	10.85	10.82	10.21	11.867	0.14	010101
					Variable	10.67	10.67	10.21	11.359	0.06	
	2	334220	1454	DROOLS	Seq	42.31	42.12	41.508	43.322	0.21	11010
					Par	11.62	11.53	10.299	14.126	0.85	2010
				PERLATO	Rule	18.86	19.02	15.794	21.549	2.07	
					Atomic	15.62	15.58	15.016	16.476	0.11	11110110
					Variable	15.08	15.02	14.24	16.689	0.21	V 1.9 AMEL 151
			572	DROOLS	Seq	19.51	19.44	18.906	21.115	0.24	10101100
					Par	6.05	6.19	3.947	8.242	1.32	10101100
	3	247113		PERLATO	Rule	8.33	8.22	7.025	9.442	0.42	1010101010
					Atomic	7.24	7.13	6.837	8.259	0.13	1414141
					Variable	6.79	6.76	6.662	7.165	0.02	01101013
	2	16 22	95	DROOLS	Seq	20.4	20.32	20.023	20.97	0.06	V11V1V1
					Par	19.52 19.71	19.52	19.115 19.232	20.072 20.657	0.07	110110101
				PERLATO	Rule	19.71	19.69 19.61	19.232	20.037	0.15	LIVI IVIN
					Atomic Variable	19.62	19.01	19.144	20.227	0.08	11110101
					Seq	42.48	42.48	41.553	43.63	0.08	11110101
				DROOLS	Par	34.27	34.6	31.866	36.019	1.13	1011010
TAXC				PERLATO	Rule	36.87	36.83	36.016	37.812	0.19	ITVI IVIVI
IAAC					Atomic	36.86	36.77	36.007	37.775	0.21	0101101
					Variable	36.8	36.74	35.972	37.698	0.23	0101103
	3	10	66	DROOLS	Seq	11.78	11.73	11.549	12.344	0.04	1303018
					Par	9.85	9.87	9.116	10.314	0.05	Jordera
				PERLATO	Rule	9.93	9.9	9.641	10.431	0.03	1342
					Atomic	9.91	9.92	9.145	10.342	0.04	Tabucdecs101/olts1
					Variable	9.98	9.97	9.309	10.761	0.06	illibordere(0)_atti

Conclusions

We created a novel solution for enhancing performance and reliability of rule-driven applications.

The results suggest that PERLATO is effective, since we achieved over 225% speedup on average.

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